No calculators or books. Submit your answers on a separate paper.

1. Consider the following algorithm. Assume A[1...n] is an array that contains integers.

```
STUPID(A, p, q)
          if p \ge q return
1
          i := p + 1
          while i < q do
3
             if A[i] < A[p] then
                SWAP(A[i], A[p])
5
             endif
6
             if A[i] > A[q] then
7
                SWAP(A[i], A[q])
8
             endif
9
             i := i + 1
10
          endwhile
11
12
          STUPID(A, p+1, q-1)
```

- (a) (2 points) Write a recurrence equation for the time consumption of Dummy(A, 1, n).
- (b) (2 points) What does the algorithm do to the array A?
- (c) (4 points) Change the algorithm so that it is no longer recursive. Give invariants to both loops in the iterative version.
- 2. Give the  $\Theta$ , or O and  $\Omega$  class of the following functions. Be as accurate as possible.
  - (a) (2 points)  $\sum_{i=1}^{n} \sqrt{i}$
  - (b) (2 points)  $\log(n!)$
  - (c) (2 points)  $\sum_{i=1}^{\lfloor \log n \rfloor} i$
  - (d) (2 points)  $O(n^2) + \Theta(n \log n)$
  - (e) (2 points)  $n^2 \log n + n^3$
- 3. Solve the following recurrences, assumin T(n) is a constant for n < 2. Simply give the answer in  $\Theta$ .
  - (a) (2 points) T(n) = T(n-1) + n/2
  - (b) (2 points) T(n) = T(n/2) + n
  - (c) (2 points) T(n) = 3T(n/3) + n
  - (d) (2 points) T(n) = 4T(n/2) + n
  - (e) (2 points)  $T(n) = T(\sqrt{n}) + \sqrt{n}$
- 4. Let (V, E) be a directed graph. We wish to represent the graph by either using adjacency matrix or adjacency list. A list requires 32 bits for each vertex and 64 bits for each edge. A matrix, on the other hand, requires only 8 bits for each slot.

- (a) (4 points) Assuming you wish to minimize the amount of memory needed for your graph, how does your choice (matrix or list) depend on the number of edges in the graph?
- (b) (4 points) We are solving a problem about graph, and the problem can be solved by an algorithm, that looks like the following. (Q is a priority queue.)

```
F(V, E)
                Q := V
1
2
                while Q \neq \emptyset do
3
                   v := Q.Extract-min()
                   for u \in V
5
                      if (u,v) \in E
6
                      endif
8
                   endfor
9
                endwhile
10
```

Assume that everything that is omitted (marked by  $\cdots$ ) is constant time, and the assignment Q:=V is  $\Theta(V)$ , and Q.Extract-min() is  $\Theta(\log n)$  when Q contains n elements, and nothing is ever added to Q. Analyse the time consumption for this algorithm for both matrices and lists.

5. In the following assignment,  $T_1$  and  $T_2$  are red-black trees and A is an array. Assume emptying a tree is constant time. T(max) () returns the largest element in the tree.

```
Dummy2(A[1,\ldots,n])
            if n \leq 1 then return
1
            T_1 := \emptyset
2
            T_2 := \emptyset
3
            T_2.insert(A[1])
4
            for k=2 to n do
5
               if A[k] < A[k \le 1] then
                  T_1.insert(A[k])
7
                   A[k]:=T_2.\mathsf{max}()
8
               else
9
                  T_2.insert(A[k])
10
               endif
11
               k := k + 1
12
            endwhile
13
```

- (a) (4 points) Analyse the time consumption of the algorithm.
- (b) (4 points) Argue, using an invariant, that after the program is run, the resulting array A is sorted.

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max.	8	10	10	8	8	44
op.		1412.0				